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Notes on the Finite Abelian HSP Algorithm

1 Introduction We quickly introduce the necessary notions, facts and the quantum algorithm, with appropriate citations. Recall the finite Abelian Hidden Subgroup Problem (HSP): Given a finite Abelian group $(G, +)$, a subgroup $H \leq G$ and some $f\colon G \to X$ with X an appropriate set, s.t. $f|_{gH}$ is constant and $f|_{gH} = f|_{hH} \to g = h$ for all $g, h \in G$. Our goal is to find a generator $\Gamma \subseteq H$ for H using a quantum algorithm[1].

(i) Since the left cosets of H induce a partition of G [2, pp. 36-37], choosing X, s.t. $|X| \geq |G/H| = |G|/|H|$ [2, p. 38], e.g. via $X := \{0,1\}^x$ for some $x \in \mathbb{N}_{\geq 1}$, $x \geq \lceil \log_2(|G/H|) \rceil$ suffices.

(ii) How do we store group elements in G in a quantum register? We can do that using qudits, because $G \cong \bigoplus_{j=1}^k \mathbb{Z}_{N_j}$ with $k \in \mathbb{N}_{\geq 1}$ and $\{N_1, ..., N_k\} \subseteq \mathbb{N}_{\geq 2}$ [2, pp. 132-135], where we take the direct sum of the groups, i.e. the elements of G can be taken to be tuples $G \ni g := (g_1, ..., g_k) \in \prod_{j=1}^k \mathbb{Z}_{N_j}$ [2, pp. 53-54]. Note that we also call $N_1, ..., N_k$ *elementary divisors*. We take such a decomposition and appropriate qudits as given here[2].

To formulate the quantum algorithm, an analogon for the $\mathbb{Z}_N$ Quantum Fourier Transform for G must be defined. This is done via characteristics.

2 Characteristics

Definition 1 ([1, p. 17])**.** A *characteristic* over G is a group homomorphism $(G, +) \to (\mathbb{C}^* := \mathbb{C} \setminus \{0\}, \cdot)$.

Lemma 1 ([1, p. 18]). The following statements are true.

(i) The *set of characteristics of* G, $\chi(G) := \{\chi\colon G \to \mathbb{C}^* \mid \chi \text{ is a characteristic over } G\}$, equipped with the composition of maps, is a group.

(ii) The map $G \hookrightarrow\!\!\!\!\!\rightarrow \chi(G), g \mapsto \chi_g$ is a group isomorphism, where we call $\chi_g\colon G \to \mathbb{C}^*, h \mapsto \prod_{j=1}^k \omega_{N_j}^{g_j h_j}$ the *characteristic induced by* g.

3 Orthogonal Subgroups

Definition 2 ([1, p. 18])**.** For $H \subseteq G$ a subgroup of a group G, we define its *orthogonal subgroup* as

$$H^\perp := \{g \in G \mid \chi_g(H) = \{1\}\} \tag{3.1}$$

Lemma 2 ([1, pp. 19-20]). The following statements hold.

(i) $H^\perp \leq G$

(ii) $H^\perp \cong G/H$

(iii) $H^{\perp\perp} = H$

Note that we included statement (i) here to justify the name in the definition.

4 General Fourier Transform

Definition 3 ([1, p. 20])**.** We define the *Quantum Fourier Transform of the Group* G as

$$\mathrm{QFT}_G := \frac{1}{|G|} \sum_{g,h \in G} \chi_g(h) \, |g\rangle\langle h| \in \mathbb{C}^{|G| \times |G|} \tag{4.1}$$

For $G = \mathbb{Z}_N$, $N \in \mathbb{N}_{\geq 1}$, we thus have $|G| = N$ and $\chi_g(h) = e^{i 2\pi \frac{gh}{N}}$ for any $g, h \in G$, meaning that this corresponds to the Quantum Fourier Transform QFT_N. We further set $|H'\rangle := \frac{1}{|H'|} \sum_{h \in H} |h\rangle$ for any subgroup $H' \leq G$. Also, we have $\mathrm{QFT}_G |0\rangle = |G|^{-1/2} \sum_{g \in G} |g\rangle$ by definition.

Lemma 3 ([1, pp. 19-21, p. 23]). The following statements are true.

(i) QFT_G is unitary.

[1] Classically, this problem is difficult, as the prime factorization problem shows [1, p. 24].

[2] Finding such a decomposition is difficult, although a quantum algorithm exists [1, p. 17].

(ii) We have $\mathrm{QFT}_G = \bigotimes_{j=1}^k \mathrm{QFT}_{\mathbb{Z}_{N_j}} = \bigotimes_{j=1}^k \mathrm{QFT}_{N_j}$ for a finite Abelian group G as in Section 1, (ii).

(iii) $\mathrm{QFT}_G \, |H\rangle = |H^\perp\rangle$

Note that in Lemma 3 (ii), each quantum fourier transform acts on a single qudit. If we only allow prime qudits, we may use the decomposition of G into cyclic groups of prime power order [2, p. 136]. Statement (iii) of Lemma 3 compactly describes the action of the general fourier group on a subgroup: It flips the group into its orthogonal complement. Applying QFT_G then again gives $|H\rangle$ by Lemma 2 (iii).

$$\begin{tikzcd} |0\rangle \arrow[r, mapsto, bend left, "H_G"] \arrow[r, mapsto, bend right, "\mathrm{QFT}_G"'] \arrow[r, phantom, "="] & |G|^{-1/2} \sum_{g\in G} |g\rangle \end{tikzcd} \qquad \begin{tikzcd} |H\rangle \arrow[r, mapsto, bend left, "\mathrm{QFT}_G"] & |H^\perp\rangle \arrow[l, mapsto, bend left, "\mathrm{QFT}_G"] \end{tikzcd}$$

In the figure, H_G denotes the Hadamard operator for G, which may be defined by the natural generalization $|h\rangle \mapsto |G|^{-1/2} \sum_{g\in G} \prod_{j=1}^k (-1)^{g_j h_j} \, |g\rangle$ for any $h \in G$. There is one more additional property that is useful.

Lemma 4 ([1, p. 20-21]). Setting for any $t \in G$

$$\tau_t := \sum_{g\in G} |t+g\rangle\langle g| \text{ and } \phi_t := \sum_{g\in G} \chi_g(t) \, |g\rangle\langle g| \tag{4.2}$$

to be its associated translation and phase shifting operators, we have the commutation

$$\mathrm{QFT}_G \tau_t = \phi_t \mathrm{QFT}_G \tag{4.3}$$

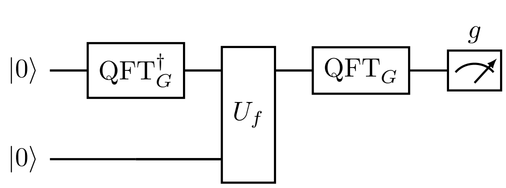
5 The Quantum Algorithm We now present the full quantum algorithm along with an analysis. The following is due to [1, pp. 22-23].

Algorithm 1 Quantum Algorithm for Solving the Finite Abelian HSP

Given: A finite Abelian group in its cyclic decomposition $G = \bigoplus_{j=1}^k \mathbb{Z}_{N_j}$ with $\{N_1, ..., N_k\} \subseteq \mathbb{N}_{\geq 2}$, $k \in \mathbb{N}_{\geq 1}$, a function $f\colon G \to X$ hiding a subgroup $H \leq G$ as described in Section 1 with $X := \{0,1\}^x$, $x \in \mathbb{N}_{\geq 1}$, $x \geq \lceil \log_2(|G|/|H|) \rceil$, a qudit register $|\Phi\rangle := |0\rangle\, |0\rangle \in S(\bigotimes_{j=1}^k \mathbb{C}^{N_j} \otimes \mathbb{C}^{|X|})$ and an oracle $U_f \in \mathbb{C}^{|G||X| \times |G||X|}$ with $|g\rangle\, |h\rangle \mapsto |g\rangle\, |h \oplus f(g)\rangle$ for all $g \in G, h \in X$.

Return: A generator $\Gamma \subseteq G$ for H.

1: $|\Phi\rangle \leftarrow (\mathrm{QFT}_G^\dagger \otimes E_{|X|})\, |\Phi\rangle$
2: $|\Phi\rangle \leftarrow U_f \, |\Phi\rangle$
3: $|\Phi\rangle \leftarrow (\mathrm{QFT}_G \otimes E_{|X|})\, |\Phi\rangle$
4: Measure $|\Phi\rangle$ wrt. the observable $\{\mathrm{Span}(\{|g\rangle\, |h\rangle \mid h \in X\} \mid g \in G)\}$ and obtain an index element $g \in G$.
5: Collect $1 + \log_2(|G|) =: t_1$ elements $g^1, ..., g^{t_1} \in G$ by repeating steps 1 to 4.
6: Form the equation system $Ah := \left(\alpha_j g_j^i\right)_{\substack{1\leq i\leq t_1 \\ 1\leq j\leq k}} (h_j)_{1\leq j\leq k} = 0$ with $h \in G$ and $\alpha_j := d/N_j$ for any $j \in \{1, ..., k\}$, where $d := \mathrm{lcm}(\{N_1, ..., N_k\})$. Compute the SNF $D \in \mathbb{Z}_d^{t_1 \times k}$ of A, and associated unimodular matrices $U \in \mathbb{Z}_d^{t_1 \times t_1}$ and $V \in \mathbb{Z}_d^{k \times k}$.
7: Sample $1 + \log_2(|G|) =: t_2$ random solutions $h^1, ..., h^{t_2}$ to the equation system $Dh' \equiv 0 \bmod d$ for $h' \in G$.
8: **return** $\{Vh^1, ..., Vh^{t_2}\}$



Note that we used the notation $S(C^n) := \{x \in \mathbb{C}^n \mid \|x\| = 1\}$ for any $n \in \mathbb{N}$.

Algorithm Analysis of the Quantum Part Let $T \subseteq G$ be a transversal wrt. G/H, i.e. a set of representatives of the induced partition. Applying the first steps yields

$$|0\rangle\,|0\rangle \xmapsto{\mathrm{QFT}_G^\dagger \otimes E_{|X|}} \frac{1}{\sqrt{|G|}} \sum_{g\in G} |g\rangle\,|0\rangle \tag{5.1}$$

$$\xmapsto{U_f} \frac{1}{\sqrt{|G|}} \sum_{g\in G} |g\rangle\,|f(g)\rangle = \frac{1}{\sqrt{|T|}} \sum_{t\in T} |t+H\rangle\,|f(t)\rangle = \frac{1}{\sqrt{|T|}} \sum_{t\in T} \tau_t\,|H\rangle\,|f(t)\rangle \tag{5.2}$$

$$\xmapsto{\mathrm{QFT}_G \otimes E_{|X|}} \frac{1}{\sqrt{|T|}} \sum_{t\in T} \mathrm{QFT}_G \tau_t\,|H\rangle\,|f(t)\rangle \overset{(1)}{=} \frac{1}{\sqrt{|H^\perp|}} \sum_{t\in T} \phi_t\,\left|H^\perp\right\rangle |f(t)\rangle \tag{5.3}$$

(1) Use the commutation relation from Lemma 4, apply Lemma 3 (iii) and then use the fact that $|T| = |G/H| = |H^\perp|$ by Lemma 2 (ii).

Note the phase shifting operator ϕ_t for any $t \in T$ in the resulting state does not influence measurements, so we have successfully, using the general QFT and the oracle, stored a uniform superposition of the elements in $\left|H^\perp\right\rangle$ in the first register. This suggests that we may repeatedly measure on this register to obtain random elements from $H^\perp$. We will apply the following lemma on random generators.

Lemma 5 ([1, pp. 76-77]). Let G be a finite group and $t \in \mathbb{N}$. Then for $t + \lceil \log_2(|G|) \rceil$ uniformly randomly chosen elements $g_1, ..., g_{t+\lceil \log_2(|G|)\rceil} \in G$, we have

$$\Pr\big(\langle g_1, ..., g_{t+\lceil \log_2(|G|)\rceil} \rangle = G\big) \geq 1 - \frac{1}{2^t} \tag{5.4}$$

Better results for this exist [1, p. 77], but this lemma suffices. However, it is still not clear how to obtain a generator for H.

Obtaining a Generator Assume for now that we have obtained elements $g^1, ..., g^\ell \in G$ with some $\ell \in \mathbb{N}_{\geq 1}$, s.t. $\langle g^1, ..., g^\ell \rangle = H^\perp$. Since $H = H^{\perp\perp}$, we have by definition for any $h \in G$, that $h \in H$, iff $\chi_h(g^j) = 1$ for any $j \in \{1, ..., \ell\}$, as annihilating a generator suffices for the definition of being in the orthogonal complement. We first reformulate the solution condition via the orthogonal complement in terms of a linear system by norming the complex roots we consider. Let $d := \mathrm{lcm}(\{N_1, ..., N_k\})$ be the least common multiple of the elementary divisors of G. Fix for now some $j' \in \{1, ..., \ell\}$. Let $\alpha_{j'} := d/N_{j'}$. Then $\omega_{N_j} = e^{i2\pi/N_j} = \omega_d^{\alpha_{j'}}$. Furthermore, $\chi_h(g^{j'}) = \prod_{j=1}^k \omega_d^{\alpha_j h_j g_j^{j'}} = 1$, iff $\sum_{j=1}^k \alpha_j h_j g_j^{j'} \equiv 0 \bmod d$. Letting j' be loose now, giving the system of congruences

$$\begin{aligned} \textstyle\sum_{j=1}^k \alpha_j g_j^1 h_j &\equiv 0 \bmod d \\ \textstyle\sum_{j=1}^k \alpha_j g_j^2 h_j &\equiv 0 \bmod d \\ &\vdots \\ \textstyle\sum_{j=1}^k \alpha_j g_j^\ell h_j &\equiv 0 \bmod d \end{aligned} \tag{5.5}$$

or in matrix notation $\left(\alpha_j g_j^i\right)_{\substack{1\leq i\leq \ell \\ 1\leq j\leq k}} (h_j)_{1\leq j\leq k} = 0 =: Ah$ over $\mathbb{Z}_d$, where we now interpret h as a column vector. If we are able to obtain enough solutions to this system of congruences, we can generate H with high probability. The necessary solution technique is the *Smith Normal Form.* Let R be a principal ideal ring and $m, n, d \in \mathbb{N}_{\geq 1}$ for the following few definitions and theorems.

Definition 4 ([3, p. 1069])**.** We define the following notions.

(i) An invertible square matrix $A \in R^{n\times n}$, $n \in \mathbb{N}_{\geq 1}$, is called *unimodular*.

(ii) Let $A \in R^{m\times n}$ be some matrix, $r := \mathrm{rk}(A)$ in the associated R-module and let $U \in R^{m\times m}$, $V \in R^{n\times n}$ be unimodular. A matrix $D \in R^{m\times n}$, s.t.

$$D = UAV = \begin{pmatrix} s_1 & & & & & \\ & \ddots & & & & \\ & & s_r & & & \\ & & & 0 & & \\ & & & & \ddots & \\ & & & & & 0 \end{pmatrix} \tag{5.6}$$

where the omitted entries in either width or height, depending on $m \le n$ or $m > n$, are zero and $s_i | s_{i+1}$ for any $i \in \{1, ..., r-1\}$, is called the *Smith Normal Form* (SNF) of A.

Theorem 6 ([3, pp. 1073-1074])**.** Any matrix $A \in \mathbb{Z}^{m\times n}$ admits to a SNF, which can be computed in time $\tilde{O}(m^3 n \log_2(m))$, additionally obtaining the similarity matrices.

With $\tilde{O}$, we denote a looser version of the O-notation, in this case omitting a few logarithmic factors. This is not the best possible result, see e.g. [4, pp. 273-274], but it is one of the simpler algorithms recovering the unimodular similarity matrices as well in the general, thus possibly singular, case.

Remark 7. In [1, p. 23], it is stated that [4] gives algorithms for computing both the SNF of a matrix over $\mathbb{Z}$, as well as the equivalence matrices, but the latter is contradicted in [4, p. 268].

Let $t_1, t_2 \in \mathbb{N}_{\ge 1}$ be loose. We first obtain $t_1 + \log_2(|G|)$ elements generating $H^\perp$ with probability $\ge 1 - 1/2^{t_1}$. After computing the SNF $D = UAV$, we have $U^{-1}DV^{-1} = A$, where we interpret first A over $\mathbb{Z}_d$ and then $\mathbb{Z}$, assuming $d \in O(1)$ in our runtime considerations. Afterwards, we interpret all matrices over $\mathbb{Z}_d$ again. We obtain a uniformly random solution to the diagonal inversion problem $Dh' \equiv 0 \bmod d$ and set $h := Vh'$, yielding $Ah = U^{-1}DV^{-1}Vh' = 0$. Thus, we may obtain $t_2 + \log_2(|G|)$ elements of H this way, which form a generator with probability $\ge 1 - 1/2^{t_2}$. In total, we obtain a generator of H with probability $\ge (1 - 1/2^{t_1})(1 - 1/2^{t_2})$. Letting $t_1 = t_2 = 1$, this is $\ge 1/4$, which means that we may execute the algorithm four times in expectation.
Runtime Analysis We apply $2k \in O(\log_2(N))$ QFT gates and use an oracle call in the first three steps of the algorithm. Note that we assume efficient implementations for the qudit unitaries $\mathrm{QFT}_{N_1}, ..., \mathrm{QFT}_{N_k}$, as these operations are local. The main cost thus stems from the computation of the SNF. Since $t_1 \in O(\log_2(|G|))$, we require a runtime of $\tilde{O}(\log_2^3(|G|)\log_2^{(2)}(|G|))$, where $\log_2^{(2)} := \log_2 \circ \log_2$.

Theorem 8. Algorithm 1 solves the HSP for a finite Abelian group with probability $\ge 1/4$ in time $\tilde{O}(\log_2^3(|G|)\log_2^{(2)}(|G|))$.

References

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